

HAMMING DISTANCE BETWEEN THE STRINGS GENERATED BY ADJACENCY MATRIX OF GENERALISED COMPLEMENTS OF A GRAPH AND THEIR SUM

SWATI NAYAK¹, SHANKAR UPADHYAY², SABITHA D'SOUZA^{1,*},
AND PRADEEP G. BHAT¹

ABSTRACT. Let $A(G)$ be the adjacency matrix of a graph G . Let $s(v)$ denote the row entries of $A(G)$ corresponding to the vertex v of G . The Hamming distance between the strings $s(u)$ and $s(v)$ is the number of positions in which $s(u)$ and $s(v)$ differ. In this paper, we study the Hamming distance between the strings generated by the adjacency matrix of generalised complements of a graph. We also compute sum of Hamming distances between all pairs of strings generated by the adjacency matrix of G_k^P and $G_{k(i)}^P$.

2010 MATHEMATICS SUBJECT CLASSIFICATION.05C99.

KEYWORDS AND PHRASES. Hamming distance, string, k -complement, $k(i)$ -complement.

1. INTRODUCTION

Let $Z_2 = \{0, 1\}$ and $(Z_2, +)$ be the additive group, where $+$ denotes addition modulo 2. For any positive integer n , $Z_2^n = Z_2 \times Z_2 \times \dots \times Z_2$ (n factors) = $\{(x_1, x_2, \dots, x_n) | x_1, x_2, \dots, x_n \in Z_2\}$. Thus every element of Z_2^n is an n -tuple (x_1, x_2, \dots, x_n) written as $x = x_1x_2\dots x_n$, where every x_i is either 0 or 1, called a string or word. The number of 1's in $x = x_1x_2\dots x_n$ is called the weight of x , denoted by $w(x)$. Let $x = x_1x_2\dots x_n$ and $y = y_1y_2\dots y_n$ be the elements of Z_2^n . The sum $x + y$ is computed by adding the corresponding components of x and y under addition modulo 2. In other words, $x_i + y_i = 0$ if $x_i = y_i$ and $x_i + y_i = 1$ if $x_i \neq y_i, i = 1, 2, \dots, n$. The Hamming distance $H_d(x, y)$ between the strings $x = x_1x_2\dots x_n$ and $y = y_1y_2\dots y_n$ is the number of i 's such that $x_i \neq y_i, 1 \leq i \leq n$. Thus $H_d(x, y) =$ Number of positions in which x and y differ = $w(x + y)$.

Example 1. Let $x = 01001$ and $y = 11010$. Therefore $x + y = 10011$. Hence $H_d(x, y) = w(x + y) = 3$.

2. PRELIMINARIES

A graph G with vertex set $V(G)$ is called a Hamming graph [2, 3] if each vertex $v \in V(G)$ can be labeled by a string $s(v)$ of a fixed length such that $H_d(s(u), s(v)) = d_G(u, v)$ for all $u, v \in V(G)$, where $d_G(u, v)$ is the length of shortest path joining u and v in G .

* corresponding author.

Submission date. 12-12-2020.

The degree of a vertex v , denoted by $deg_G(v)$ is the number of neighbours of v also denoted by $n(v)$. The vertices which are adjacent to both u and v are called common neighbours of u and v . The vertices which are neither adjacent to u nor adjacent to v are called non common neighbours of u and v . We denote two vertices u adjacent to v by $u \sim v$ and u not adjacent to v by $u \not\sim v$.

Sum of Hamming distances between all pairs of strings generated by the adjacency matrix of a graph G is denoted by $HA(G)$ is given by,

$$HA(G) = \sum_{1 \leq i < j \leq n} H_d(s(v_i), s(v_j)).$$

$HA(G)$ is a graph invariant.

Definition 2.1. [5] For a graph $G(V, E)$, let $P = \{V_1, V_2, \dots, V_k\}$ be the partition of vertex set $V(G)$ of order $k \geq 2$. For each set V_r in P , remove the edges of graph G inside V_r and add the edges of the complement of G joining the vertices of V_r . The graph thus obtained is called the $k(i)$ -complement of graph G with respect to partition P denoted by $G_{k(i)}^P$.

Definition 2.2. [4] Let $G(V, E)$ be any graph. Let $P = \{V_1, V_2, \dots, V_k\}$ be a partition of V of order $k \geq 2$. For all V_i and V_j in $P, i \neq j$, remove the edges between V_i and V_j in G and add the edges between V_i and V_j which are not in G . The graph thus obtained is called k -complement of G with respect to P denoted by G_k^P .

Throughout the paper, $H_{d_{k(i)}}(s(u), s(v))$ and $H_{d_k}(s(u), s(v))$ denote Hamming distance between the strings $s(u)$ and $s(v)$ generated by adjacency matrix of $G_{k(i)}^P$ and G_k^P respectively.

3. HAMMING DISTANCE BETWEEN THE STRINGS IN $k(i)$ -COMPLEMENT AND k -COMPLEMENT OF A GRAPH.

Theorem 3.1. Let $G(V, E)$ be any graph of order n and $S \subseteq V$. Let $P = \{S, V - S\}$ be the partition of vertex set $V(G)$.

- (1) Suppose the vertices u and v have $k_1(k_2)$ common neighbours, $l_1(l_2)$ non-common neighbours respectively in $\langle S \rangle (\langle V - S \rangle)$.
 - (a) If $u \sim v$ in G and $u, v \in S$, then $H_{d_{2(i)}}(s(u), s(v)) = n - (k_1 + k_2) - (l_1 + l_2) - 2$.
 - (b) If $u \not\sim v$ in G and $u, v \in S$, then $H_{d_{2(i)}}(s(u), s(v)) = n - (k_1 + k_2) - (l_1 + l_2)$.
- (2) Suppose $r_1 = n(u), r_2 = n(v)$ and $r_3 = n(u), r_4 = n(v)$ in $\langle S \rangle$ and $\langle V - S \rangle$ respectively.
 - (a) If $u \sim v$ in G and $u \in S, v \in V - S$, then $H_{d_{2(i)}}(s(u), s(v)) = n - (r_1 + r_2 + r_3 + r_4)$.
 - (b) If $u \not\sim v$ in G and $u \in S, v \in V - S$, then $H_{d_{2(i)}}(s(u), s(v)) = n - (r_1 + r_2 + r_3 + r_4) - 2$.

Proof. Let G be any graph of order n and $S \subseteq V$. Let $P = \{S, V - S\}$ be the partition of vertex set V .

- (1) Let the vertices u, v of G be have k_1 common neighbours, l_1 non common neighbours in $\langle S \rangle$ and k_2 common neighbours, l_2 non common neighbours in $\langle V - S \rangle$.

(a) If $u \sim v$ in G and $u, v \in S$, then $u \approx v$ in $G_{2(i)}^P$. We observe that u and v have $(k_2 + l_1)$ common neighbours and $(k_1 + l_2)$ non common neighbours. Remaining $n - (k_2 + l_1) - (k_1 + l_2) - 2$ vertices other than u and v are adjacent to either u or v . Therefore strings of u and v from $A(G_{2(i)}^P)$ will be of the form

$$s(u) = x_1x_2x_3 \dots x_{k_2+l_1+1}x_{k_2+l_1+2}x_{k_2+l_1+3} \dots x_{k_2+l_1+k_1+l_2+2}x_{k_2+l_1+k_1+l_2+3} \dots x_n \text{ and}$$

$$s(v) = y_1y_2y_3 \dots y_{k_2+l_1+1}y_{k_2+l_1+2}y_{k_2+l_1+3} \dots y_{k_2+l_1+k_1+l_2+2}y_{k_2+l_1+k_1+l_2+3} \dots y_n,$$

where $x_1 = 0, x_2 = 0, y_1 = y_2 = 0, x_i = y_i = 1$ for $i = 3, 4, \dots, k_2 + l_1 + 2, x_i = y_i = 0$ for $i = k_2 + l_1 + 3, k_2 + l_1 + 4, \dots, k_2 + l_1 + k_1 + l_2 + 2$ and $x_i \neq y_i$ for $i = k_2 + l_1 + k_1 + l_2 + 3, k_2 + l_1 + k_1 + l_2 + 4, \dots, n$. Thus $s(u)$ and $s(v)$ differ at $n - (k_2 + l_2) - (k_1 + l_2) - 2$ places.

$$\text{Hence } H_{d_{2(i)}}(s(u), s(v)) = n - (k_1 + k_2) - (l_1 + l_2) - 2.$$

(b) If $u \approx v$ in G and $u, v \in S$, then $u \sim v$ in $G_{2(i)}^P$. The vertices u and v have $(k_2 + l_1)$ common neighbours and $(k_1 + l_2)$ non common neighbours. Remaining $n - (k_2 + l_1) - (k_1 + l_2) - 2$ vertices are adjacent to either u or v . Therefore strings of u and v from $A(G_{2(i)}^P)$ will be of the form

$$s(u) = x_1x_2x_3 \dots x_{k_2+l_1+1}x_{k_2+l_1+2}x_{k_2+l_1+3} \dots x_{k_2+l_1+k_1+l_2+2}x_{k_2+l_1+k_1+l_2+3} \dots x_n.$$

$$\text{and } s(v) = y_1y_2y_3 \dots y_{k_2+l_1+1}y_{k_2+l_1+2}y_{k_2+l_1+3} \dots y_{k_2+l_1+k_1+l_2+2}y_{k_2+l_1+k_1+l_2+3} \dots y_n.$$

Where $x_1 = 0, x_2 = 1, y_1 = 1, y_2 = 0, x_i = y_i = 1$ for $i = 3, 4, \dots, k_2 + l_1 + 2, x_i = y_i = 0$ for $i = k_2 + l_1 + 3, k_2 + l_1 + 4, \dots, k_2 + l_1 + k_1 + l_2 + 2$ and $x_i \neq y_i$ for $i = k_2 + l_1 + k_1 + l_2 + 3, k_2 + l_1 + k_1 + l_2 + 4, \dots, n$. Thus $s(u)$ and $s(v)$ differ at $n - (k_2 + l_2) - (k_1 + l_2) - 2 + 2$ places. Hence $H_{d_{2(i)}}(s(u), s(v)) = n - (k_1 + k_2) - (l_1 + l_2)$.

- (2) Let $r_1 = n(u), r_2 = n(v)$ and $r_3 = n(u), r_4 = n(v)$ in $\langle S \rangle$ and $\langle V - S \rangle$ respectively.

(a) If $u \sim v$ in G and $u \in S, v \in V - S$, then $u \sim v$ in $G_{2(i)}^P$. The vertices u and v have $(r_2 + r_3)$ common neighbours and $(r_1 + r_4)$ non common neighbours. Remaining $n - (r_2 + r_3) - (r_1 + r_4) - 2$ vertices are adjacent to either u or v . Therefore strings of u and v from $A(G_{2(i)}^P)$ will be of the form

$$s(u) = x_1x_2x_3 \dots x_{r_2+r_3+1}x_{r_2+r_3+2}x_{r_2+r_3+3} \dots x_{r_2+r_3+r_1+r_4+1}x_{r_2+r_3+r_1+r_4+2}x_{r_2+r_3+r_1+r_4+3} \dots x_n \text{ and}$$

$$s(v) = y_1y_2y_3 \dots y_{r_2+r_3+1}y_{r_2+r_3+2}y_{r_2+r_3+3} \dots y_{r_2+r_3+r_1+r_4+1}$$

$y_{r_2+r_3+r_1+r_4+2}y_{r_2+r_3+r_1+r_4+3} \dots y_n$, where $x_1 = 0, x_2 = 1, y_1 = 1, y_2 = 0, x_i = y_i = 1$ for $i = 3, 4, \dots, (r_2 + r_3 + 2)$ and $x_i = y_i = 0$ for $i = (r_2 + r_3 + 3), (r_2 + r_3 + 4), \dots, (r_2 + r_3 + r_1 + r_4 + 2), x_i \neq y_i$

for $i = (r_2 + r_3 + r_1 + r_4 + 3), \dots, n$. Thus $s(u)$ and $s(v)$ differ at $n - (r_2 + r_3) - (r_1 + r_4) - 2 + 2$ places.

Hence $H_{d_2(i)}(s(u), s(v)) = n - (r_1 + r_2 + r_3 + r_4)$.

- (b) If $u \approx v$ in G and $u \in S, v \in V - S$, then $u \approx v$ in $G_{2(i)}^P$. The vertices u and v have $(r_2 + r_3)$ common neighbours and $(r_1 + r_4)$ non common neighbours. Remaining $n - (r_2 + r_3) - (r_1 + r_4) - 2$ vertices are adjacent to either u or v . Therefore strings of u and v from $A(G_{2(i)}^P)$ will be of the form

$$s(u) = x_1x_2x_3 \dots x_{r_2+r_3+1}x_{r_2+r_3+2}x_{r_2+r_3+3} \dots x_{r_2+r_3+r_1+r_4+1}x_{r_2+r_3+r_1+r_4+2}x_{r_2+r_3+r_1+r_4+3} \dots x_n \text{ and}$$

$$s(v) = y_1y_2y_3 \dots y_{r_2+r_3+1}y_{r_2+r_3+2}y_{r_2+r_3+3} \dots y_{r_2+r_3+r_1+r_4+1}y_{r_2+r_3+r_1+r_4+2}y_{r_2+r_3+r_1+r_4+3} \dots y_n.$$

Where $x_1 = 0, x_2 = 0, y_1 = 0, y_2 = 0, x_i = y_i = 1$ for $i = 3, 4, \dots, (r_2 + r_3 + 2)$ and $x_i = y_i = 0$ for $i = (r_2 + r_3 + 3), (r_2 + r_3 + 4), \dots, (r_2 + r_3 + r_1 + r_4 + 2), x_i \neq y_i$ for $i = (r_2 + r_3 + r_1 + r_4 + 3), \dots, n$. Thus $s(u)$ and $s(v)$ differ at $n - (r_2 + r_3) - (r_1 + r_4) - 2$ places.

Hence $H_{d_2(i)}(s(u), s(v)) = n - (r_1 + r_2 + r_3 + r_4) - 2$.

□

Corollary 3.2. Let $G(V, E)$ be any graph of order n and $S \subseteq V$. Let $P = \{S, V - S\}$ be the partition of vertex set $V(G)$.

- (1) Suppose the vertices u, v of G have k_1 common neighbours, l_1 non common neighbours in $\langle S \rangle$ and k_2 common neighbours, l_2 non common neighbours in $\langle V - S \rangle$.
 - (a) If $u \sim v$ in G and $u, v \in S$, then $H_{d_2}(s(u), s(v)) = n - (k_1 + k_2) - (l_1 + l_2)$.
 - (b) If $u \approx v$ in G and $u, v \in S$, then $H_{d_2}(s(u), s(v)) = n - (k_1 + k_2) - (l_1 + l_2) - 2$.
- (2) Suppose $r_1 = n(u), r_2 = n(v)$ in $\langle S \rangle$ and $r_3 = n(u), r_4 = n(v)$ in $\langle V - S \rangle$.
 - (a) If $u \sim v$ in G and $u \in S, v \in V - S$, then $H_{d_2}(s(u), s(v)) = n - (r_1 + r_2 + r_3 + r_4) - 2$.
 - (b) If $u \approx v$ in G and $u \in S, v \in V - S$, then $H_{d_2}(s(u), s(v)) = n - (r_1 + r_2 + r_3 + r_4)$.

Theorem 3.3. Let $G(V, E)$ be any graph of order n and $S \subseteq V$. Let $P = \{S, V - S\}$ be the partition of vertex set $V(G)$.

- (1) Let the vertices u, v of G have k_1 common neighbours, l_1 non common neighbours in $\langle S \rangle$ and k_2 common neighbours, l_2 non common neighbours in $\langle V - S \rangle$.
 - (a) If $u \sim v$ in G and $u, v \in S$, then $deg_{G_{2(i)}^P}(u) + deg_{G_{2(i)}^P}(v) = n + (k_2 + l_1) - (k_1 + l_2) - 2$.
 - (b) If $u \approx v$ in G and $u, v \in S$, then $deg_{G_{2(i)}^P}(u) + deg_{G_{2(i)}^P}(v) = n + (k_2 + l_1) - (k_1 + l_2)$.
- (2) Let $r_1 = n(u), r_2 = n(v)$ in $\langle S \rangle$ and $r_3 = n(u), r_4 = n(v)$ in $\langle V - S \rangle$.

- (a) If $u \sim v$ in G and $u \in S, v \in V - S$, then $deg_{G_{2(i)}^P}(u) + deg_{G_{2(i)}^P}(v) = n + r_2 + r_3 - (r_1 + r_4)$.
- (b) If $u \not\sim v$ in G and $u \in S, v \in V - S$ in $\langle S \rangle$, then $deg_{G_{2(i)}^P}(u) + deg_{G_{2(i)}^P}(v) = n + r_2 + r_3 - (r_1 + r_4) - 2$.

Proof. Let $G(V, E)$ be any graph of order n and $S \subseteq V$. Let $P = \{S, V - S\}$ be the partition of vertex set $V(G)$.

- (1) Let the vertices u, v of G have k_1 common neighbours, l_1 non common neighbours in $\langle S \rangle$ and k_2 common neighbours, l_2 non common neighbours in $\langle V - S \rangle$.
 - (a) If $u \sim v$ in G and $u, v \in S$, then $u \approx v$ in $G_{2(i)}^P$. The vertices u and v have $(k_2 + l_1)$ common neighbours and $(k_1 + l_2)$ non common neighbours. Remaining $n - (k_2 + l_1) - (k_1 + l_2) - 2$ vertices are adjacent to either u or v . Therefore

$$deg_{G_{2(i)}^P}(u) + deg_{G_{2(i)}^P}(v) = (k_2 + l_1) + (k_2 + l_1) + n - (k_2 + l_1) - (k_1 + l_2) - 2 = n + (k_2 + l_1) - (k_1 + l_2) - 2.$$
 - (b) If $u \not\sim v$ in G and $u, v \in S$, then $u \sim v$ in $G_{2(i)}^P$. The vertices u and v have $(k_2 + l_1)$ common neighbours and $(k_1 + l_2)$ non common neighbours. Remaining $n - (k_2 + l_1) - (k_1 + l_2) - 2$ vertices are adjacent to either u or v . Therefore

$$deg_{G_{2(i)}^P}(u) + deg_{G_{2(i)}^P}(v) = (K_2 + l_1) + (k_2 + l_1) + n - (k_2 + l_1) - (k_1 + l_2) - 2 + 2 = n + (k_2 + l_1) - (k_1 + l_2).$$
- (2) Let $r_1 = n(u), r_2 = n(v)$ in $\langle S \rangle$ and $r_3 = n(u), r_4 = n(v)$ in $\langle V - S \rangle$.
 - (a) If $u \sim v$ in G and $u \in S, v \in V - S$, then $u \sim v$ in $G_{2(i)}^P$. The vertices u and v have $(r_2 + r_3)$ common neighbours and $(r_1 + r_4)$ non common neighbours. Remaining $n - (r_2 + r_3) - (r_1 + r_4) - 2$ vertices are adjacent to either u or v . Therefore

$$deg_{G_{2(i)}^P}(u) + deg_{G_{2(i)}^P}(v) = (r_2 + r_3) + (r_2 + r_3) + n - (r_2 + r_3) - (r_1 + r_4) - 2 + 2 = n + r_2 + r_3 - (r_1 + r_4).$$
 - (b) If $u \not\sim v$ in G and $u \in S, v \in V - S$, then $u \not\sim v$ in $G_{2(i)}^P$. The vertices u and v have $(r_2 + r_3)$ common neighbours and $(r_1 + r_4)$ non common neighbours. Remaining $n - (r_2 + r_3) - (r_1 + r_4) - 2$ vertices are adjacent to either u or v . Therefore

$$deg_{G_{2(i)}^P}(u) + deg_{G_{2(i)}^P}(v) = (r_2 + r_3) + (r_2 + r_3) + n - (r_2 + r_3) - (r_1 + r_4) - 2 = n + r_2 + r_3 - (r_1 + r_4) - 2.$$

□

Corollary 3.4. Let $G(V, E)$ be any graph of order n and $S \subseteq V$. Let $P = \{S, V - S\}$ be the partition of vertex set $V(G)$.

- (1) Suppose the vertices u, v of G have k_1 common neighbours, l_1 non common neighbours in $\langle S \rangle$ and k_2 common neighbours, l_2 non common neighbours in $\langle V - S \rangle$.
 - (a) If $u \sim v$ in G and $u, v \in S$, then $deg_{G_2^P}(u) + deg_{G_2^P}(v) = n + (k_1 + l_2) - (k_2 + l_1)$.
 - (b) If $u \not\sim v$ in G and $u, v \in S$, then $deg_{G_2^P}(u) + deg_{G_2^P}(v) = n + (k_1 + l_2) - (k_2 + l_1) - 2$.

- (2) Suppose $r_1 = n(u), r_2 = n(v)$ in $\langle S \rangle$ and $r_3 = n(u), r_4 = n(v)$ in $\langle V - S \rangle$.
 - (a) If $u \sim v$ in G and $u \in S, v \in V - S$, then $deg_{G_2^P}(u) + deg_{G_2^P}(v) = n + r_1 + r_4 - (r_2 + r_3) - 2$.
 - (b) If $u \not\sim v$ in G and $u \in S, v \in V - S$, then $deg_{G_2^P}(u) + deg_{G_2^P}(v) = n + r_1 + r_4 - (r_2 + r_3) - 2$.

Theorem 3.5. For every vertices u and v in graph G ,

$$|H_{d_k}(s(u), s(v)) - H_{d_{k(i)}}(s(u), s(v))| = 2.$$

Proof. We consider two cases here.

Case 1: Let $u \sim v$ in G_k^P . Then $u \not\sim v$ in $G_{k(i)}^P$. We know that $\overline{G_k^P} \cong G_{k(i)}^P$ and $\overline{G_{k(i)}^P} \cong G_k^P$ [5]. If u and v have k_1 common and l_1 non common neighbours in G_k^P , then u and v have l_1 common and k_1 non common neighbours in $G_{k(i)}^P$. So $H_{d_k}(s(u), s(v)) = n - k_1 - l_1$ and $H_{d_{k(i)}}(s(u), s(v)) = n - l_1 - k_1 - 2$. Therefore

$$(1) \quad H_{d_k}(s(u), s(v)) - H_{d_{k(i)}}(s(u), s(v)) = 2.$$

Case 2: Let $u \not\sim v$ in G_k^P . Then $u \sim v$ in $G_{k(i)}^P$. If u and v have k_2 common and l_2 non common neighbours in G_k^P , then u and v have l_2 common and k_2 non common neighbours in $G_{k(i)}^P$. So $H_{d_k}(s(u), s(v)) = n - k_2 - l_2 - 2$ and $H_{d_{k(i)}}(s(u), s(v)) = n - l_2 - k_2$. Therefore

$$(2) \quad H_{d_{k(i)}}(s(u), s(v)) - H_{d_k}(s(u), s(v)) = 2.$$

Hence from equations (1) and (2), we get

$$|H_{d_k}(s(u), s(v)) - H_{d_{k(i)}}(s(u), s(v))| = 2.$$

□

Theorem 3.6. For a graph G on n vertices,

$$HA(G_{k(i)}^P) = HA(G_k^P) + n(n - 1) - 4q.$$

Where q denotes the number of edges of G_k^P .

Proof. We have $\overline{G_k^P} \cong G_{k(i)}^P$ and $\overline{G_{k(i)}^P} \cong G_k^P$ [5].

Also $HA(\overline{G}) = HA(G) + n(n - 1) - 4q$ [1].

Therefore $HA(G_{k(i)}^P) = HA(G_k^P) + n(n - 1) - 4q$. □

4. SUM OF HAMMING DISTANCES OF 2 AND 2(i) COMPLEMENT OF SOME GRAPHS

Theorem 4.1. For a complete graph K_n , $HA((K_n)_2^P) = (n - m)(mn - 4m) + n(n - 1)$.

Proof. Let $P = \{V_1, V_2\}$ be a partition of vertex set of K_n , where $V_1 = \{v_1, v_2, \dots, v_m\}$ and $V_2 = \{v_m, v_{m+1}, \dots, v_n\}$.

We observe that $(K_n)_2^P = K_m \cup K_{n-m}$.

Let us consider following cases.

Case 1: Let $u, v \in V_1$.

Then by Theorem (4) of [1],

$$\sum_{u,v \in V_1} H_{d_2}(s(u), s(v)) = m(m - 1).$$

Case 2: Let $u, v \in V_2$. By Theorem (4) of [1],

$$\sum_{u,v \in V_2} H_{d_2}(s(u), s(v)) = (n - m)(n - m - 1).$$

Case 3: Let $u \in V_1$ and $v \in V_2$. In $(K_n)_2^P$, $u \approx v$. There are $m(n - m)$ pairs of vertices with zero common and zero non common neighbours which have Hamming distance $n - 2$.

Therefore

$$\begin{aligned} HA((K_n)_2^P) &= \sum_{u,v \in V_1} H_{d_2}(s(u), s(v)) + \sum_{u,v \in V_2} H_{d_2}(s(u), s(v)) \\ &\quad + \sum_{u \in V_1, v \in V_2} H_{d_2}(s(u), s(v)) \\ &= m(m - 1) + (n - m)(n - m - 1) + m(n - m)(n - 2) \\ &= (n - m)(mn - 4m) + n(n - 1). \end{aligned}$$

□

Corollary 4.2. $HA((K_n)_{2(i)}^P) = mn(n - m)$.

Theorem 4.3. *If $P = \{V_1, V_2\}$ is a partition of cycle graph C_n with $\langle V_1 \rangle = P_m$ and $\langle V_2 \rangle = P_{n-m}$, then $HA((C_n)_2^P) = (n - m)(mn - 8m) + 2(n^2 - 2n + 8)$.*

Proof. Let $V_1 = \{v_1, v_2, \dots, v_m\}$ and $V_2 = \{v_{m+1}, v_{m+2}, \dots, v_n\}$ be the partition of C_n .

Case 1: Let $u, v \in V_1$. Assume that $u \sim v$ in C_n . Then $u \sim v$ in $(C_n)_2^P$. There are $m - 1$ pairs of vertices in C_n with distance 1. Out of $m - 1$ pairs of vertices, 2 pairs (v_1, v_2) and (v_{m-1}, v_m) have $n - m - 1$ common neighbours and $m - 3$ non common neighbours in $(C_n)_2^P$. Hence $H_{d_2}(s(u), s(v)) = n - (n - m - 1) - (m - 3) = 4$.

Remaining $m - 3$ pairs of vertices have $n - m$ common and $m - 4$ non common neighbours in $(C_n)_2^P$. Thus they have Hamming distance 4. Therefore

$$\sum_{d_{C_n}=1} H_{d_2}(s(u), s(v)) = 4m - 4.$$

Let $d_{C_n}(u, v) = 2$, where $u \approx v$ in C_n and $(C_n)_2^P$. There are $m - 2$ pairs of vertices which are at distance 2. Out of which 2 pairs (v_1, v_3) and (v_{m-2}, v_m) have $n - m$ common neighbours and $m - 4$ non common neighbours in $(C_n)_2^P$. So $H_{d_2}(s(u), s(v)) = n - (n - m) - (m - 4) - 2 = 2$.

Remaining $m - 4$ pairs of vertices have $n - m + 1$ common neighbours and $m - 5$ non common neighbours. Hence $H_{d_2}(s(u), s(v)) = n - (n - m + 1) - (m - 5) - 2 = 2$.

Therefore

$$\sum_{d_{C_n}=2} H_{d_2}(s(u), s(v)) = 2m - 4.$$

Let $d_{C_n}(u, v) > 2$. We observe that $u \approx v$ in $(C_n)_2^P$. There are $\binom{m}{2} - 2m + 3$ pairs of vertices which are at distance greater than 2. Out of these vertices, $2(m - 4)$ pairs of vertices have Hamming distance 4 as they have $n - m - 1$ common and $m - 5$ non common neighbours and one pair (v_1, v_m) have $n - m - 2$ common and $m - 4$ non common neighbours which results in $H_{d_2}(u, v) = 4$. Remaining $\binom{m}{2} - 4m + 10$ pairs of vertices have $n - m$ common and $m - 6$ non common neighbours. So $H_{d_2}(u, v) = 4$. Hence

$$\sum_{d_{C_n} > 2} H_{d_2}(s(u), s(v)) = 2m(m - 1) - 8m + 12.$$

Therefore

$$\begin{aligned} \sum_{u, v \in V_1} H_{d_2}(s(u), s(v)) &= \sum_{d_{C_n}=1} H_{d_2}(s(u), s(v)) + \sum_{d_{C_n}=2} H_{d_2}(s(u), s(v)) \\ &+ \sum_{d_{C_n} > 2} H_{d_2}(s(u), s(v)) \\ &= 4m - 4 + 2m - 4 + 2m(m - 1) - 8m + 12 \\ &= 2m^2 - 4m + 4. \end{aligned}$$

Case 2: Let $u, v \in V_2$. Then

$$\sum_{u, v \in V_2} H_{d_2}(s(u), s(v)) = 2(n - m)^2 - 4(n - m) + 4.$$

Case 3: Let $u \in V_1$ and $v \in V_2$ such that $u \sim v$ in C_n . Then $u \approx v$ in $(C_n)_2^P$. There are 2 pairs of vertices (v_1, v_n) and (v_m, v_{m+1}) of distance 1 in C_n . Since these two pairs of vertices have 2 common and zero non common neighbours in $(C_n)_2^P$, Hamming distance is $n - 4$.

Let $u \approx v$ in C_n . Then $u \sim v$ in $(C_n)_2^P$. There are $m(n - m) - 2$ such pairs of vertices. Among these pairs, (v_1, v_{m+1}) and (v_m, v_n) have 2 common and 2 non common neighbours. Hence these two pairs have Hamming distance $n - 4$. The four pairs of vertices (v_m, v_{m+2}) , (v_1, v_{n-1}) , (v_n, v_2) and (v_{m+1}, v_{m-1}) have Hamming distance $n - 2$ as they have 2 common and zero non common neighbours. There are $2n - 12$ pairs of vertices having three common and one non common neighbours, which are of the form (v_1, v_j) , $j = m + 2, m + 3, \dots, n - 2$, (v_m, v_j) , $j = m + 3, m + 4, \dots, n - 1$, (v_i, v_{m+1}) , $i = 2, 3, \dots, m - 2$ and (v_i, v_{n-m}) , $i = 3, 4, \dots, m - 1$ with Hamming distance $n - 4$. Remaining $m(n - m) - 2n + 4$ pairs of vertices have 4 common and zero non common neighbours and thus have Hamming distance $n - 4$. So,

$$\sum_{u \in V_1, v \in V_2} H_{d_2}(s(u), s(v)) = (m(n - m) + 4n - 24)(n - 4) + 4(n - 2).$$

Therefore,

$$\begin{aligned}
 HA((C_n)_2^P) &= \sum_{u,v \in V_1} H_{d_2}(s(u), s(v)) + \sum_{u,v \in V_2} H_{d_2}(s(u), s(v)) + \\
 &\quad \sum_{u \in V_1, v \in V_2} H_{d_2}(s(u), s(v)) \\
 &= 2m^2 - 4m + 4 + 2(n - m)^2 - 4(n - m) + 4 + (m(n - m) + 4n - 24) \\
 &\quad (n - 4) + 4(n - 2) \\
 &= (n - m)(mn - 8m) + 2(n^2 - 2n + 8).
 \end{aligned}$$

□

Corollary 4.4. *If $P = \{V_1, V_2\}$ is a partition of cycle graph C_n with $\langle V_1 \rangle = P_m$ and $\langle V_2 \rangle = P_{n-m}$, then $HA((C_n)_2^P) = (n - m)(mn - 12m) + 3n^2 - 9n + 32$.*

Theorem 4.5. *Let $P = \{V_1, V_2\}$ be a partition of path graph P_n such that $V_1 = \langle P_m \rangle$ and $V_2 = \langle P_{n-m} \rangle$. Then $HA((P_n)_2^P) = (n - m)(mn - 8m) + 2(n^2 - 2n + 7)$.*

Proof. Let $P = \{V_1, V_2\}$ be a partition of path graph P_n such that $V_1 = \{v_1, v_2, \dots, v_m\}$ and $V_2 = \{v_{m+1}, v_{m+2}, \dots, v_n\}$.

We consider following cases.

Case 1: Let $u, v \in V_1$, $u \sim v$ in P_n . They remain adjacent in $(P_n)_2^P$.

There are $m - 1$ pairs of vertices of distance 1. Out of $m - 1$ pairs, a pair of vertices (v_{m-1}, v_m) has $n - m - 1$ common neighbours and $m - 3$ non common neighbours. Thus Hamming distance of (v_{m-1}, v_m) is 4. Also (v_1, v_2) have $n - m$ common neighbours and $m - 3$ non common neighbours. Hence Hamming distance of (v_1, v_2) is 3. Remaining $m - 3$ pairs of vertices have $n - m$ common and $m - 4$ non common neighbours. Hence, Hamming distance of these pairs of vertices is 4. Therefore

$$\sum_{d_{v_n}=1} H_{d_2}(s(u), s(v)) = 4m - 5.$$

There are $m - 2$ pairs of vertices which are at distance 2. These vertices are non adjacent in $(P_n)_2^P$. Out of $m - 2$ pairs of vertices, (v_{m-2}, v_m) has $n - m$ common and $m - 4$ non common neighbours and (v_1, v_3) has $n - m + 1$ common and $m - 4$ non common neighbours. So Hamming distance of (v_{m-2}, v_m) is 2 and that of (v_1, v_3) is 1. Remaining $m - 4$ pairs of vertices have $n - m + 1$ common and $m - 5$ non common neighbours. Hence Hamming distance of such pairs of vertices is 2.

$$\sum_{d_{v_n}=2} H_{d_2}(s(u), s(v)) = 2m - 5.$$

There are $\binom{m}{2} - 2m + 3$ pairs of vertices which are at distance greater than 2. Since $m - 4$ pairs of vertices of the form (v_1, v_i) , $i = 4, 5, \dots, m - 1$ have $n - m$ common and $m - 5$ non common neighbours, Hamming distance of (v_1, v_i) is 3. Similarly, the Hamming distance of pairs of vertices of the form (v_j, v_m) , $j = 2, 3, \dots, m - 3$ is 4, since (v_j, v_m) pairs have $n - m - 1$ common

and $m - 5$ non common neighbours. Hamming distance of (v_1, v_m) is 4 as it has $n - m - 1$ common and $m - 4$ non common neighbours. Remaining $\binom{m}{2} - 4m + 10$ pairs of vertices have $n - m$ common and $m - 6$ non common neighbours. Hence their Hamming distance is 4.

$$\sum_{d_{p_n} > 2} H_{d_2}(s(u), s(v)) = 2m^2 - 11m + 15.$$

Therefore

$$\begin{aligned} \sum_{u, v \in V_1} H_{d_2}(s(u), s(v)) &= \sum_{d_{p_n} = 1} H_{d_2}(s(u), s(v)) + \sum_{d_{p_n} = 2} H_{d_2}(s(u), s(v)) \\ &\quad + \sum_{d_{p_n} > 2} H_{d_2}(s(u), s(v)) \\ &= 2m^2 - 5m + 5. \end{aligned}$$

Case 2: If $u, v \in V_2$, then

$$\sum_{u, v \in V_2} H_{d_2}(s(u), s(v)) = 2(n - m)^2 - 5(n - m) + 5.$$

Case 3: Let $u \in V_1$ and $v \in V_2$. There are $m(n - m)$ pairs of vertices between V_1 and V_2 . Among these pairs, consider (v_m, v_{m+1}) and (v_1, v_n) . Now $v_m \sim v_{m+1}$ and $v_1 \approx v_n$ in P_n . Also $v_m \approx v_{m+1}$ and $v_1 \sim v_n$ in $(P_n)_2^P$. So both the pairs have 2 common and zero non common neighbours. Thus they have Hamming distance $n - 4$ and $n - 2$ respectively. Also (v_m, v_n) , (v_1, v_{m+1}) have 2 common and one non common neighbours and thus have Hamming distance $n - 3$. The pairs (v_{m-1}, v_{m+1}) , (v_m, v_{m+2}) have Hamming distance $n - 2$ as they have 2 common and zero non common neighbours. Further $n - 4$ pairs of vertices which are of the form (v_1, v_i) , $i = m + 2, m + 3, \dots, n - m - 1$ and (v_j, v_n) , $j = 2, 3, \dots, m - 1$ have Hamming distance $n - 4$ as they have 3 common and zero non common neighbours. There are $n - 6$ pairs of vertices which are of the form (v_i, v_{m+1}) , $i = 2, 3, \dots, m - 2$ and (v_m, v_j) , $j = m + 3, m + 4, \dots, n - m - 1$, having 3 common and 1 non common neighbours. Hence they have Hamming distance $n - 4$. Remaining $m(n - m) - 2n + 4$ pairs of vertices have 4 common and zero non common neighbours. Hence Hamming distance of these pairs of vertices is $n - 4$. Thus

$$\begin{aligned} \sum_{u \in V_1, v \in V_2} H_{d_2}(s(u), s(v)) &= (n - 4) + (n - 4)(n - 3) + 2(n - 3) + 3(n - 2) \\ &\quad + (n - 6)(n - 4) \\ &\quad + (m(n - m) - 2n + 4)(n - 4) \\ &= mn(n - m) + 4m^2 - 4mn + n + 4. \end{aligned}$$

Therefore

$$\begin{aligned}
 HA((P_n)_2^P) &= \sum_{u,v \in V_1} H_{d_2}(s(u), s(v)) + \sum_{u,v \in V_2} H_{d_2}(s(u), s(v)) \\
 &\quad + \sum_{u \in V_1, v \in V_2} H_{d_2}(s(u), s(v)) \\
 &= 2m^2 - 5m + 5 + 2(n - m)^2 - 5(n - m) + 5 + mn(n - m) \\
 &\quad + 4m^2 - 4mn + n + 4 \\
 &= (n - m)(mn - 8m) + 2(n^2 - 2n + 7).
 \end{aligned}$$

□

Corollary 4.6. *Let $P = \{V_1, V_2\}$ be a partition of path graph P_n such that $V_1 = \langle P_m \rangle$ and $V_2 = \langle P_{n-m} \rangle$. Then $HA((P_n)_{2(i)}^P) = (n - m)(mn - 12m) + 3n^2 - 9n + 26$.*

Theorem 4.7. *Let $P = \{V_1, V_2\}$ be a partition of complete bipartite graph $K_{r,s}$ such that $\langle V_1 \rangle = K_{a,b}$ and $\langle V_2 \rangle = K_{r-a,s-b}$. Then $HA((K_{r,s})_2^P) = (r + s)[(r - s)(b - a) - (a - b)^2 + rs]$.*

Proof. Let $\langle V_1 \rangle = K_{a,b}$ and $\langle V_2 \rangle = K_{r-a,s-b}$.

Case 1: Let $u, v \in V_1$. Suppose $u \sim v$ in $K_{r,s}$. As there are ab pairs of vertices in $(K_{r,s})_2^P$ having no common as well as non common neighbours, the Hamming distance between u and v is $r + s$.

If $u \not\sim v$ in $K_{r,s}$, then we see that 2-complement of $K_{r,s}$ will have $\binom{a}{2}$ pairs of vertices with $b + r - a$ common and $s - b + a - 2$ non common neighbours and $\binom{b}{2}$ pairs of vertices with $a + s - b$ common and $r - a + b - 2$ non common neighbours. Thus $H_{d_2}(s(u), s(v)) = 0$ for every (u, v) in $(K_{r,s})_2^P$. Hence

$$\begin{aligned}
 \sum_{u,v \in V_1} H_{d_2}(s(u), s(v)) &= ab(r + s) + \binom{a}{2}(0) + \binom{b}{2}(0) \\
 &= ab(r + s).
 \end{aligned}$$

Case 2: Let $u, v \in V_2$. Then $\sum_{u,v \in V_2} H_{d_2}(s(u), s(v)) = (r - a)(s - b)(r + s)$.

Case 3: Let $u \in V_1$ and $v \in V_2$. Then there are $a(r - a) + b(s - a)$ pairs of vertices with Hamming distance $r + s$ as they have zero common and zero non common neighbours and remaining pairs of vertices have Hamming distance zero. So

$$\sum_{u \in V_1, v \in V_2} H_{d_2}(s(u), s(v)) = (r + s)(a(r - a) + b(s - a)).$$

Therefore

$$\begin{aligned} HA((K_{r,s})_2^P) &= \sum_{u,v \in V_1} H_{d_2}(s(u), s(v)) + \sum_{u,v \in V_2} H_{d_2}(s(u), s(v)) \\ &\quad + \sum_{u \in V_1, v \in V_2} H_{d_2}(s(u), s(v)) \\ &= ab(r+s) + (r-a)(s-b)(r+s) + (r+s)(a(r-a) + b(s-a)) \\ &= (r+s)[(s-r)(b-a) - (a-b)^2 + rs]. \end{aligned}$$

□

Corollary 4.8. *Let $P = \{V_1, V_2\}$ be a partition of complete bipartite graph $K_{r,s}$ such that $\langle V_1 \rangle = K_{a,b}$ and $\langle V_2 \rangle = K_{r-a, s-b}$. Then $HA((K_{r,s})_{2(i)}^P) = b(4r - r^2 + a(2r - 8)) - r - a(-r^2 + 4r) + s^2(b - a + r + 1) - s(2r - 4(a - b) + (a - b)^2 - r^2 + 1) + r^2 - (a^2 + b^2)(r - 4)$.*

Corollary 4.9. *Let $P = \{V_1, V_2\}$ be a partition of complete bipartite graph $K_{r,s}$ with $\langle V_1 \rangle = rK_1$ and $\langle V_2 \rangle = sK_1$. Then $HA((K_{r,s})_2^P) = 0$.*

Proof. Since $(K_{r,s})_2^P$ is totally disconnected graph with respect to P ,

$$HA((K_{r,s})_2^P) = 0. \quad \square$$

Corollary 4.10. *Let $P = \{V_1, V_2\}$ be a partition of complete bipartite graph $K_{r,s}$ with $\langle V_1 \rangle = rK_1$ and $\langle V_2 \rangle = sK_1$. Then $HA((K_{r,s})_{2(i)}^P) = (r+s)(r+s-1)$.*

REFERENCES

- [1] A. B. Ganagi and H. S. Ramane, *Hamming distance between the strings generated by adjacency matrix of a graph and their sum*, Algebra Discrete Math. 22 (2016), 82-93.
- [2] W. Imrich and S. Klavžar, *A simple $O(mn)$ algorithm for recognizing Hamming graphs*, Bull. Inst. Combin. Appl. 9 (1993), 45-56.
- [3] S. Klavžar and I. Peterin, *Characterizing subgraphs of Hamming graphs*, J. Graph Theory, 49 (2005), 302-312.
- [4] E. Sampathkumar and L. Pushpalatha, *Complement of a graph: A Generalization*, Graphs Combin. 14 (1998), 377-392.
- [5] E. Sampathkumar, L. Pushpalatha, C. V. Venkatachalam and P. G. Bhat, *Generalized complements of a graph*, Indian J. Pure Appl. Math. 29 (1998), 625-639.

1. DEPARTMENT OF MATHEMATICS, MANIPAL INSTITUTE OF TECHNOLOGY, MANIPAL ACADEMY OF HIGHER EDUCATION, MANIPAL, INDIA.

Email address: swati.nayak@manipal.edu, sabitha.dsouza@manipal.edu, pg.bhat@manipal.edu

2. DEPARTMENT OF MATHEMATICS, MAHARAJA INSTITUTE OF TECHNOLOGY, MYSORE

Email address: shankar.upadhyay8@gmail.com